

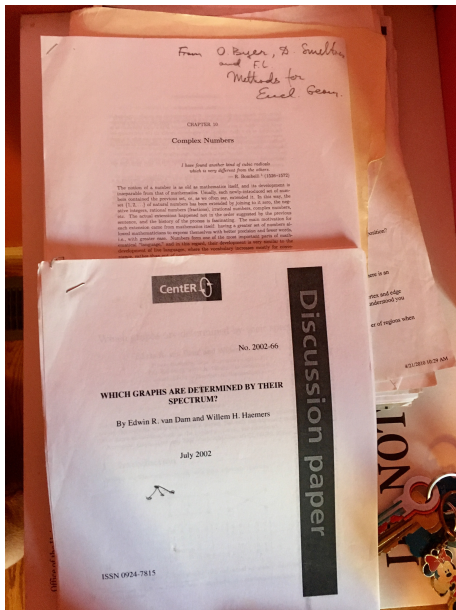
# Degree Ramsey numbers for even cycles

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**AEGT: a conference in honor of Willem, Felix, and Andy**  
University of Delaware

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“Large” implies “structure”

# Ramsey Theory

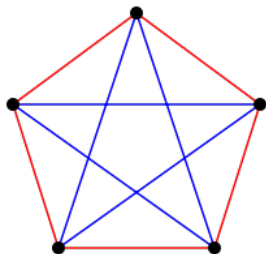
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For graphs  $G$  and  $H$ ,  $R(G, H)$  denotes the smallest  $n$  such that any red/blue coloring of  $E(K_n)$  contains a red  $G$  or a blue  $H$

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$$R(C_n, C_m) = f(n, m)$$

$$R(K_n, T_m) = (m - 1)(n - 1) + 1$$

$$R(K_3, K_t) = \Theta\left(\frac{t^2}{\log t}\right)$$

$$R(K_m, K_n) = ??$$

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- $R(G, G) = R_{|V(H)|}(G, 2)$ .
- $\rho = e(H)$ : *size Ramsey number*
- $\rho = \chi(H)$ : *chromatic Ramsey number*
- $\rho = \Delta(H)$ : *degree Ramsey number*

# Degree Ramsey numbers

- Trees, cycles (Kinnersley, Milans, West)
- $R_{\Delta}(C_4, s) = \Omega(s^{14/9})$  (Jiang, Milans, West)
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$R_{\Delta}(C_4, s) = \Omega(s^2)$ : **any** graph of maximum degree  $\Delta$  can be partitioned into  $O(\Delta^{1/2})$  subgraphs each of which is  $C_4$  free.

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Theorem (Tait)

$$R_\Delta(C_6, s) = \Theta(s^{3/2})$$

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If one can partition  $K_\Delta$  into  $C_6$  or  $C_{10}$  free graphs efficiently then one can partition any graph of maximum degree  $\Delta$  into  $C_6$  or  $C_{10}$  free graphs efficiently (this is not quite true).



Assume  $H$  is a graph of maximum degree  $\Delta$ . We would like to partition it into “few” graphs all of which are  $C_{2k}$ -free.

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### Lemma

*There is a subgraph spanning  $H'$  of  $H$  and a proper coloring of  $V(H')$  with  $O(\Delta)$  colors such that*

- 1 *For all  $v$ ,  $d_{H'}(v) \geq \delta d_H(v)$*
- 2 *Every vertex in  $H'$  sees a rainbow*

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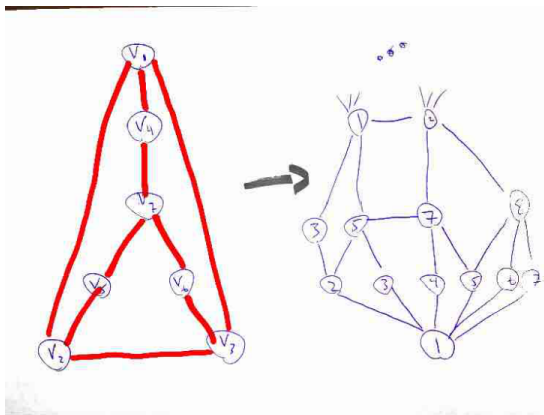
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Note: by iterating, it suffices to show that  $H'$  can be partitioned into “few” graphs all of which are  $C_{2k}$ -free.

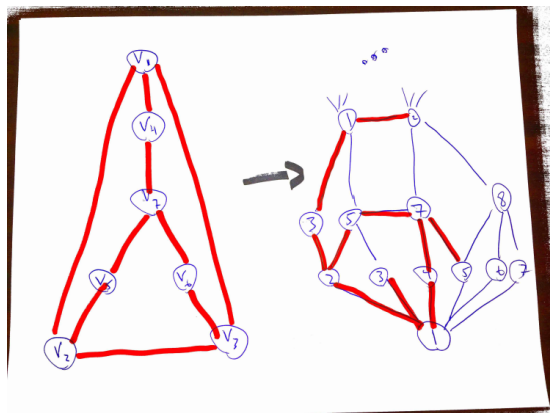
Assume  $H'$  is properly colored with  $100\Delta$  colors so that each neighborhood is a rainbow. Also assume that we have partitioned  $K_{100\Delta}$  into subgraphs  $G_1, \dots, G_m$  none of which contain  $C_4$ .

We show that we can **also partition**  $H'$  into  $m$  subgraphs none of which contain a  $C_4$ .

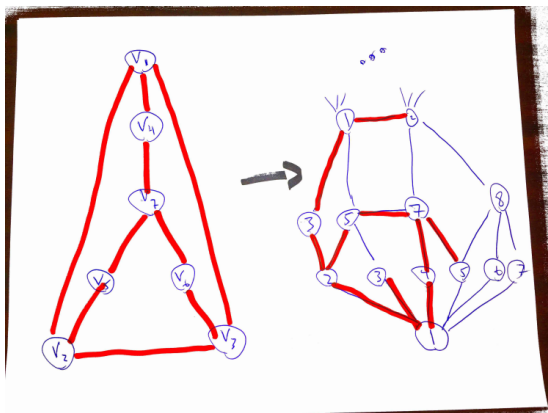
# Map from $K_\Delta$ to $H'$



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Since  $G_1, \dots, G_m$  partition  $K_{100\Delta}$ , these graphs partition  $H'$ .

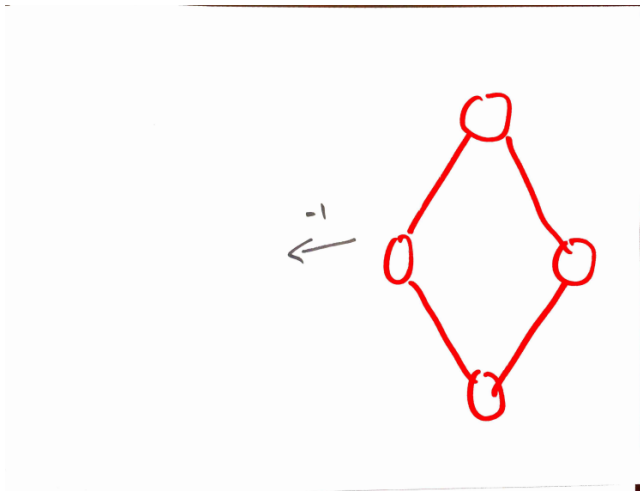
## $C_4$ free?

What is the preimage of a  $C_4$ ?



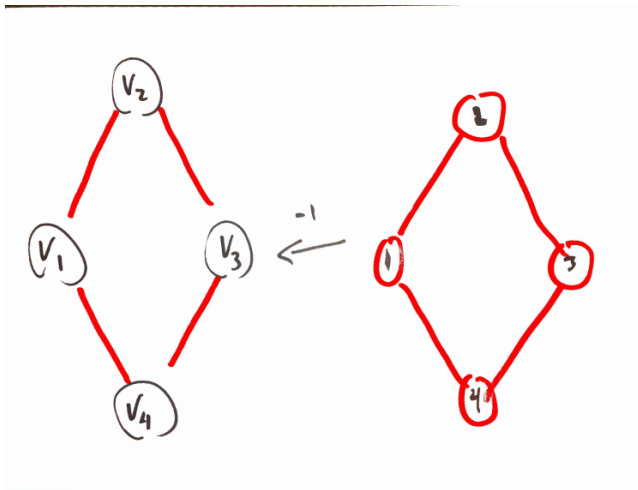
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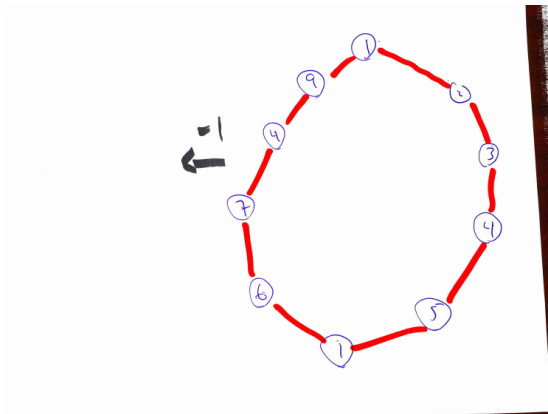


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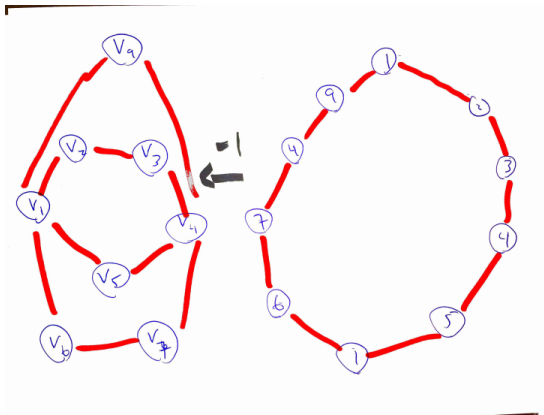
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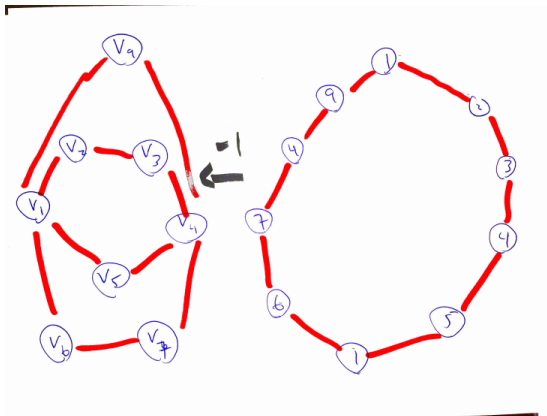
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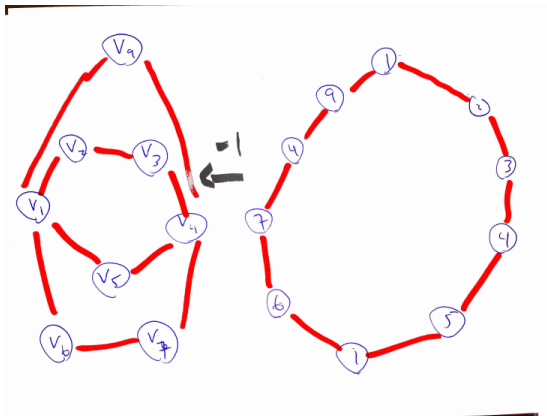
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The preimage is a **nonbacktracking closed walk** of length  $2k$ . It contains a cycle of length at most  $2k$ .

If we can partition  $K_{C\Delta}$  into “few” graphs with girth greater than  $2k$ , then we can partition  $H'$  into the same number of graphs with no  $C_{2k}$ .

Theorem (Chung-Graham 1975, Lazebnik-Woldar 2000)

*$K_n$  can be partitioned into  $m$   $C_4$  free subgraphs where  $m \sim n^{1/2}$ .*



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Theorem (Li-Lih 2009)

*$K_n$  can be partitioned into  $O(n^{2/3})$   $C_6$ -free graphs or  $O(n^{4/5})$   $C_{10}$ -free graphs.*

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Theorem

*$K_n$  can be partitioned in  $O(n^{2/3})$  subgraphs of girth 8 or  $O(n^{4/5})$  subgraphs of girth 12.*

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Theorem (Lazebnik, Ustimenko, Woldar 1995)

*Let  $k$  be fixed and  $\delta = 0$  if  $k$  is odd and 1 if  $k$  is even. The graphs  $CD(k, q)$  are graphs on  $n$  vertices with  $\Omega\left(n^{1+\frac{2}{3k-3+\delta}}\right)$  edges and girth at least  $2k + 2$ .*

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**Corollary**

Let  $k \geq 2$  and  $\delta = 0$  if  $k$  is odd and  $\delta = 1$  if  $k$  is even. Then

$$R_{\Delta}(C_{2k}, s) = \Omega\left(\left(\frac{s}{\log s}\right)^{1+\frac{2}{3k-5+\delta}}\right).$$

### Open Problem 1

Get rid of the log in the previous theorem.

### Open Problem 2

Can any graph with spectral radius  $\lambda$  be decomposed into  $O(\lambda^{2/3})$  subgraphs each of which are  $C_4$ -free?

### Open Problem 3

Is there a function  $f(\Delta, s)$  such that  $R_\Delta(G, s) \leq f(\Delta(G), s)$ ?

