Toughness, connectivity and the spectrum of regular graphs

Xiaofeng Gu (University of West Georgia)

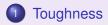
Joint work with S.M. Cioabă (University of Delaware)

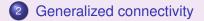
AEGT, August 7, 2017

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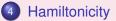
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Matrix and Eigenvalue

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- $\lambda_i(G)$ denotes the *i*th largest eigenvalue of *G*. So we have $\lambda_1 \ge \lambda_2 \ge \cdots \ge \lambda_n$.
- Let $\lambda = \max\{|\lambda_2|, |\lambda_3|, \cdots, |\lambda_n|\} = \max\{|\lambda_2|, |\lambda_n|\}.$

- - The toughness t(G) of a connected graph G is defined as $t(G) = \min\{\frac{|S|}{c(G-S)}\}$, where the minimum is taken over all proper subset $S \subset V(G)$ such that c(G - S) > 1.

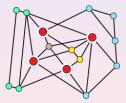


Figure: toughness = 1

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- It was disproved by Bauer, Broersma and Veldman (2000).
- Conjecture (Chvátal, 1973) There exists some positive t_0 such that any graph with toughness greater than t_0 is Hamiltonian.

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Some results

• Theorem (Alon 1995)

For any connected *d*-regular graph G, $t(G) > \frac{1}{3}(\frac{d^2}{d\lambda+\lambda^2}-1)$.

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- Conjecture (Brouwer, 1995) For any connected *d*-regular graph G, $t(G) > \frac{d}{\lambda} - 1$.

More results

Theorem (Cioabă and G. 2016)
 For any connected *d*-regular graph *G* with *d* ≥ 3 and edge connectivity κ' < d, t(G) > ^d/_{λ2} - 1 ≥ ^d/_λ - 1.

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- brief idea:

1. Let *G* be a connected *d*-regular graph with edge connectivity κ' . Then $t(G) \ge \kappa'/d$.

2. Let *G* be a *d*-regular graph with $d \ge 2$ and edge connectivity $\kappa' < d$. Then $\lambda_2(G) \ge d - \frac{2\kappa'}{d+1}$.

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• Brouwer's conjecture remains unsolved for the case $\kappa' = d$.

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More results

• Theorem (Liu and Chen 2010) For any connected *d*-regular graph *G*, if

$$\lambda_2(G) < \begin{cases} d-1+\frac{3}{d+1}, & \text{if } d \text{ is even,} \\ d-1+\frac{2}{d+1}, & \text{if } d \text{ is odd,} \end{cases}$$

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• Theorem (Cioabă and Wong 2014) For any connected *d*-regular graph *G*, if

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 Theorem (Cioabă and G. 2016) Let G be a connected d-regular graph with d ≥ 3 and edge connectivity κ'. If κ' = d, or, if κ' < d and

$$\lambda_{\lceil \frac{d}{d-\kappa'}\rceil}(G) < \begin{cases} \frac{d-2+\sqrt{d^2+12}}{2}, & \text{if } d \text{ is even,} \\ \frac{d-2+\sqrt{d^2+8}}{2}, & \text{if } d \text{ is odd,} \end{cases}$$

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then $t(G) \ge 1$.

• Theorem (Cioabă and G. 2016) For any bipartite connected *d*-regular graph *G* with $\kappa' < d$, if $\lambda_{\lceil \frac{d}{d-\kappa'}\rceil}(G) < d - \frac{d-1}{2d}$, then t(G) = 1.

Useful tools: Interlacing Theorem

• Theorem

Let *A* be a real symmetric $n \times n$ matrix and *B* be a principal $m \times m$ submatrix of *A*. Then $\lambda_i(A) \ge \lambda_i(B) \ge \lambda_{n-m+i}(A)$ for $1 \le i \le m$.

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Corollary

Let S_1, S_2, \dots, S_k be disjoint subsets of V(G) with $e(S_i, S_j) = 0$ for $i \neq j$. Then

$$\lambda_k(G) \ge \lambda_k(G[\cup_{i=1}^k S_i]) \ge \min_{1 \le i \le k} \{\lambda_1(G[S_i])\}.$$

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Generalized connectivity

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- Given an integer $l \ge 2$, Chartrand, Kapoor, Lesniak and Lick defined the *l*-connectivity $\kappa_l(G)$ of a graph *G* to be the minimum number of vertices of *G* whose removal produces a disconnected graph with at least *l* components or a graph with fewer than *l* vertices.

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- Given an integer $l \ge 2$, Chartrand, Kapoor, Lesniak and Lick defined the *l*-connectivity $\kappa_l(G)$ of a graph G to be the minimum number of vertices of G whose removal produces a disconnected graph with at least l components or a graph with fewer than l vertices.
- By definition, for a noncomplete connected graph G, we have $t(G) = \min_{2 \le l \le \alpha} \{\frac{\kappa_l(G)}{l}\}$ where α is the independence number of G.

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Results

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Results

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- Theorem (Krivelevich and Sudakov 2006) For a *d*-regular graph, $\kappa \ge d - \frac{36\lambda^2}{d}$.
- Theorem (Cioabă and G. 2016)
 Let *l*, *k* be integers with *l* ≥ *k* ≥ 2. For any connected *d*-regular graph *G* with |V(G)| ≥ *k* + *l* − 1, *d* ≥ 3 and edge connectivity κ', if κ' = *d*, or, if κ' < *d* and

$$\lambda_{\lceil \frac{(l-k+1)d}{d-\kappa'}\rceil}(G) < \begin{cases} \frac{d-2+\sqrt{d^2+12}}{2}, & \text{if } d \text{ is even} \\ \frac{d-2+\sqrt{d^2+8}}{2}, & \text{if } d \text{ is odd}, \end{cases}$$

then $\kappa_l(G) \ge k$.

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Corollaries

• Corollary (Cioabă and G. 2016) Let $l \ge 2$. For any connected *d*-regular graph *G* with $|V(G)| \ge l + 1$ and $d \ge 3$, if

$$\lambda_l(G) < \begin{cases} \frac{d-2+\sqrt{d^2+12}}{2}, & \text{if } d \text{ is even,} \\ \frac{d-2+\sqrt{d^2+8}}{2}, & \text{if } d \text{ is odd,} \end{cases}$$

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Corollary (Cioabă and G. 2016)
 For any connected *d*-regular graph *G* with *d* ≥ 3, if

$$\lambda_2(G) < \begin{cases} \frac{d-2+\sqrt{d^2+12}}{2}, & \text{if } d \text{ is even}, \\ \frac{d-2+\sqrt{d^2+8}}{2}, & \text{if } d \text{ is odd}, \end{cases}$$

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Spanning tree with bounded maximum degree

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- Theorem (Win 1989) Let $k \ge 2$ and G be a connected graph. If for any $S \subseteq V(G)$, $c(G - S) \le (k - 2)|S| + 2$, then G has a spanning k-tree.

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Spanning tree with bounded maximum degree

• Theorem (Wong 2013) Let $k \ge 3$ and G be a connected d-regular graph. If $\lambda_4 < d - \frac{d}{(k-2)(d+1)}$, then G has a spanning k-tree.

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Theorem (Wong 2013)
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 Theorem (Cioabă and G. 2016)

Let $k \geq 3$ and G be a connected d-regular graph with edge connectivity κ' . Let $l = d - (k - 2)\kappa'$. Each of the following statements holds.

(i) If $l \le 0$, then *G* has a spanning *k*-tree. (ii) If l > 0 and $\lambda_{\lceil \frac{3d}{l} \rceil} < d - \frac{d}{(k-2)(d+1)}$, then *G* has a spanning *k*-tree.

Hamiltonian graphs

• Conjecture (Krivelevich and Sudakov, 2002) Let *G* be a *d*-regular graph with *n* vertics and with the second largest absolute value λ . There exist a positive consitant *C* such that for large enough *n*, if $d/\lambda > C$, then *G* is Hamiltonian.

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- Recall: Theorem (Brouwer, 1995) For any connected *d*-regular graph G, $t(G) > \frac{d}{\lambda} - 2$.

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- Recall: Conjecture (Chvátal, 1973) There exists some positive t_0 such that any graph with toughness greater than t_0 is Hamiltonian.
- Recall: Theorem (Brouwer, 1995) For any connected *d*-regular graph G, $t(G) > \frac{d}{\lambda} - 2$.
- Krivelevich and Sudakov proved, if $d/\lambda > f(n)$, then G is Hamiltonian.

Thank You